# The CAST-256 Encryption Algorithm

Carlisle Adams

CAST-256 is a symmetric cipher designed in accordance with the CAST design procedure as outlined in [A97]. It is an extension of the CAST-128 cipher and has been submitted as a candidate for NIST's Advanced Encryption Standard (AES) effort -- see http://csrc.nist.gov/encryption/aes/aes_home.htm for details.

This document contains several sections of the CAST-256 AES Submission Package delivered to NIST on June 9th, 1998. All complete submissions received by NIST will be made public in late August at the First AES Candidate Conference, but the following material is being made available now so that public analysis of the CAST-256 algorithm may begin (see, for example, http://www.ii.uib.no/~larsr/aes.html for the current status of submitted algorithms).

Many thanks are due to those who worked with me in the (long, challenging, frustrating, and very enjoyable!) design and analysis phases that ultimately led to the detailed specification given below: Howard Heys (Memorial University); Stafford Tavares (Queen's University); and Michael Wiener (Entrust Technologies). As well, many thanks are due to the two who did the various implementations on a variety of platforms (Reference C, Optimized C, Optimized Java, and even M6811 Assembler): Serge Mister and Ian Clysdale (both of Entrust Technologies).

[A97] C. Adams, "Constructing Symmetric Ciphers Using the CAST Design Procedure", in *Selected Areas in Cryptography*, Kluwer Academic Publishers, 1997, pp.71-104 (reprinted from *Designs, Codes and Cryptography*, vol. 12, no. 3, November 1997).

# CAST-256
# Algorithm Specification

## 1. Algorithm Specification

### 1.1 CAST-128 Notation

The following notation from CAST-128 [A97b, A97c] is relevant to CAST-256.

- CAST-128 uses a pair of subkeys per round: a 5-bit quantity $k_{r_i}$ is used as a "rotation" key for round *i* and a 32-bit quantity $k_{m_i}$ is used as a "masking" key for round *i*.

- Three different round functions are used in CAST-128. The rounds are as follows (where $D$ is the data input to the operation, $I_a$ - $I_d$ are the most significant byte through least significant byte of $I$, respectively, $S_i$ is the $i^{\text{th}}$ s-box (see following page for s-box definitions), and $O$ is the output of the operation). Note that $+$ and $-$ are addition and subtraction modulo $2^{32}$, $\oplus$ is bitwise eXclusive-OR, and $\hookleftarrow$ is the circular left-shift operation.

  Type 1:

  $$I = ((k_{m_i} + D) \hookleftarrow k_{r_i})$$
  $$O = ((S_1[I_a] \oplus S_2[I_b]) - S_3[I_c]) + S_4[I_d]$$

  Type 2:

  $$I = ((k_{m_i} \oplus D) \hookleftarrow k_{r_i})$$
  $$O = ((S_1[I_a] - S_2[I_b]) + S_3[I_c]) \oplus S_4[I_d]$$

  Type 3:

  $$I = ((k_{m_i} - D) \hookleftarrow k_{r_i})$$
  $$O = ((S_1[I_a] + S_2[I_b]) \oplus S_3[I_c]) - S_4[I_d]$$

  Let $f_1, f_2, f_3$ be keyed round function operations of Types 1, 2, and 3 (respectively) above.

## *CAST-128 Notation (cont’d)*

- CAST-128 uses four round function substitution boxes (s-boxes), $S_1$ - $S_4$. These are defined as follows (entries (written in hexadecimal notation) are to be read left-to-right, top-to-bottom).

### S-Box $S_1$

```
30fb40d4 9fa0ff0b 6beccd2f 3f258c7a 1e213f2f 9c004dd3 6003e540 cf9fc949
bfd4af27 88bbbdb5 e2034090 98d09675 6e63a0e0 15c361d2 c2e7661d 22d4ff8e
28683b6f c07fd059 ff2379c8 775f50e2 43c340d3 df2f8656 887ca41a a2d2bd2d
a1c9e0d6 346c4819 61b76d87 22540f2f 2abe32e1 aa54166b 22568e3a a2d341d0
66db40c8 a784392f 004dff2f 2db9d2de 97943fac 4a97c1d8 527644b7 b5f437a7
b82cbaef d751d159 6ff7f0ed 5a097a1f 827b68d0 90ecf52e 22b0c054 bc8e5935
4b6d2f7f 50bb64a2 d2664910 bee5812d b7332290 e93b159f b48ee411 4bff345d
fd45c240 ad31973f c4f6d02e 55fc8165 d5b1caad a1ac2dae a2d4b76d c19b0c50
882240f2 0c6e4f38 a4e4bfd7 4f5ba272 564c1d2f c59c5319 b949e354 b04669fe
b1b6ab8a c71358dd 6385c545 110f935d 57538ad5 6a390493 e63d37e0 2a54f6b3
3a787d5f 6276a0b5 19a6fcdf 7a42206a 29f9d4d5 f61b1891 bb72275e aa508167
38901091 c6b505eb 84c7cb8c 2ad75a0f 874a1427 a2d1936b 2ad286af aa56d291
d7894360 425c750d 93b39e26 187184c9 6c00b32d 73e2bb14 a0bebc3c 54623779
64459eab 3f328b82 7718cf82 59a2cea6 04ee002e 89fe78e6 3fab0950 325ff6c2
81383f05 6963c5c8 76cb5ad6 d49974c9 ca180dcf 380782d5 c7fa5cf6 8ac31511
35e79e13 47da91d0 f40f9086 a7e2419e 31366241 051ef495 aa573b04 4a805d8d
548300d0 00322a3c bf64cddf ba57a68e 75c6372b 50afd341 a7c13275 915a0bf5
6b54bfab 2b0b1426 ab4cc9d7 449ccd82 f7fbf265 ab85c5f3 1b55db94 aad4e324
cfa4bd3f 2deaa3e2 9e204d02 c8bd25ac eadf55b3 d5bd9e98 e31231b2 2ad5ad6c
954329de adbe4528 d8710f69 aa51c90f aa786bf6 22513f1e aa51a79b 2ad344cc
7b5a41f0 d37cfbad 1b069505 41ece491 b4c332e6 032268d4 c9600acc ce387e6d
bf6bb16c 6a70fb78 0d03d9c9 d4df39de e01063da 4736f464 5ad328d8 b347cc96
75bb0fc3 98511bfb 4ffbcc35 b58bcf6a e11f0abc bfc5fe4a a70aec10 ac39570a
3f04442f 6188b153 e0397a2e 5727cb79 9ceb418f 1cacd68d 2ad37c96 0175cb9d
c69dff09 c75b65f0 d9db40d8 ec0e7779 4744ead4 b11c3274 dd24cb9e 7e1c54bd
f01144f9 d2240eb1 9675b3fd a3ac3755 d47c27af 51c85f4d 56907596 a5bb15e6
580304f0 ca042cf1 011a37ea 8dbfaadb 35ba3e4a 3526ffa0 c37b4d09 bc306ed9
98a52666 5648f725 ff5e569d 0ced63d0 7c63b2cf 700b45e1 d5ea50f1 85a92872
af1fbda7 d4234870 a7870bf3 2d3b4d79 42e04198 0cd0ede7 26470db8 f881814c
474d6ad7 7c0c5e5c d1231959 381b7298 f5d2f4db ab838653 6e2f1e23 83719c9e
bd91e046 9a56456e dc39200c 20c8c571 962bda1c e1e696ff b141ab08 7cca89b9
1a69e783 02cc4843 a2f7c579 429ef47d 427b169c 5ac9f049 dd8f0f00 5c8165bf
```

### S-Box $S_2$

```
1f201094 ef0ba75b 69e3cf7e 393f4380 fe61cf7a eec5207a 55889c94 72fc0651
ada7ef79 4e1d7235 d55a63ce de0436ba 99c430ef 5f0c0794 18dcdb7d a1d6eff3
a0b52f7b 59e83605 ee15b094 e9ffd909 dc440086 ef944459 ba83ccb3 e0c3cdfb
d1da4181 3b092ab1 f997f1c1 a5e6cf7b 01420ddb e4e7ef5b 25a1ff41 e180f806
1fc41080 179bee7a d37ac6a9 fe5830a4 98de8b7f 77e83f4e 79929269 24fa9f7b
e113c85b acc40083 d7503525 f7ea615f 62143154 0d554b63 5d681121 c866c359
3d63cf73 cee234c0 d4d87e87 5c672b21 071f6181 39f7627f 361e3084 e4eb573b
602f64a4 d63acd9c 1bbc4635 9e81032d 2701f50c 99847ab4 a0e3df79 ba6cf38c
10843094 2537a95e f46f6ffe a1ff3b1f 208cfb6a 8f458c74 d9e0a227 4ec73a34
fc884f69 3e4de8df ef0e0088 3559648d 8a45388c 1d804366 721d9bfd a58684bb
e8256333 844e8212 128d8098 fed33fb4 ce280ae1 27e19ba5 d5a6c252 e49754bd
c5d655dd eb667064 77840b4d a1b6a801 84db26a9 e0b56714 21f043b7 e5d05860
54f03084 066ff472 a31aa153 dadc4755 b5625dbf 68561be6 83ca6b94 2d6ed23b
eccf01db a6d3d0ba b6803d5c af77a709 33b4a34c 397bc8d6 5ee22b95 5f0e5304
81ed6f61 20e74364 b45e1378 de18639b 881ca122 b96726d1 8049a7e8 22b7da7b
5e552d25 5272d237 79d2951c c60d894c 488cb402 1ba4fe5b a4b09f6b 1ca815cf
a20c3005 8871df63 b9de2fcb 0cc6c9e9 0beeff53 e3214517 b4542835 9f63293c
ee41e729 6e1d2d7c 50045286 1e6685f3 f33401c6 30a22c95 31a70850 60930f13
```

```
73f98417 a1269859 ec645c44 52c877a9 cdff33a6 a02b1741 7cbad9a2 2180036f
50d99c08 cb3f4861 c26bd765 64a3f6ab 80342676 25a75e7b e4e6d1fc 20c710e6
cdf0b680 17844d3b 31eef84d 7e0824e4 2ccb49eb 846a3bae 8ff77888 ee5d60f6
7af75673 2fdd5cdb a11631c1 30f66f43 b3faec54 157fd7fa ef8579cc d152de58
db2ffd5e 8f32ce19 306af97a 02f03ef8 99319ad5 c242fa0f a7e3ebb0 c68e4906
b8da230c 80823028 dcdef3c8 d35fb171 088a1bc8 bec0c560 61a3c9e8 bca8f54d
c72feffa 22822e99 82c570b4 d8d94e89 8b1c34bc 301e16e6 273be979 b0ffeaa6
61d9b8c6 00b24869 b7ffce3f 08dc283b 43daf65a f7e19798 7619b72f 8f1c9ba4
dc8637a0 16a7d3b1 9fc393b7 a7136eeb c6bcc63e 1a513742 ef6828bc 520365d6
2d6a77ab 3527ed4b 821fd216 095c6e2e db92f2fb 5eea29cb 145892f5 91584f7f
5483697b 2667a8cc 85196048 8c4bacea 833860d4 0d23e0f9 6c387e8a 0ae6d249
b284600c d835731d dcb1c647 ac4c56ea 3ebd81b3 230eabb0 6438bc87 f0b5b1fa
8f5ea2b3 fc184642 0a036b7a 4fb089bd 649da589 a345415e 5c038323 3e5d3bb9
43d79572 7e6dd07c 06dfdf1e 6c6cc4ef 7160a539 73bfbe70 83877605 4523ecf1
```

S-Box $S_3$

```
8defc240 25fa5d9f eb903dbf e810c907 47607fff 369fe44b 8c1fc644 aececa90
beb1f9bf eefbcaea e8cf1950 51df07ae 920e8806 f0ad0548 e13c8d83 927010d5
11107d9f 07647db9 b2e3e4d4 3d4f285e b9afa820 fade82e0 a067268b 8272792e
553fb2c0 489ae22b d4ef9794 125e3fbc 21fffcee 825b1bfd 9255c5ed 1257a240
4e1a8302 bae07fff 528246e7 8e57140e 3373f7bf 8c9f8188 a6fc4ee8 c982b5a5
a8c01db7 579fc264 67094f31 f2bd3f5f 40fff7c1 1fb78dfc 8e6bd2c1 437be59b
99b03dbf b5dbc64b 638dc0e6 55819d99 a197c81c 4a012d6e c5884a28 ccc36f71
b843c213 6c0743f1 8309893c 0feddd5f 2f7fe850 d7c07f7e 02507fbf 5afb9a04
a747d2d0 1651192e af70bf3e 58c31380 5f98302e 727cc3c4 0a0fb402 0f7fef82
8c96fdad 5d2c2aae 8ee99a49 50da88b8 8427f4a0 1eac5790 796fb449 8252dc15
efbd7d9b a672597d ada840d8 45f54504 fa5d7403 e83ec305 4f91751a 925669c2
23efe941 a903f12e 60270df2 0276e4b6 94fd6574 927985b2 8276dbcb 02778176
f8af918d 4e48f79e 8f616ddf e29d840e 842f7d83 340ce5c8 96bbb682 93b4b148
ef303cab 984faf28 779faf9b 92dc560d 224d1e20 8437aa88 7d29dc96 2756d3dc
8b907cee b51fd240 e7c07ce3 e566b4a1 c3e9615e 3cf8209d 6094d1e3 cd9ca341
5c76460e 00ea983b d4d67881 fd47572c f76cedd9 bda8229c 127dadaa 438a074e
1f97c090 081bdb8a 93a07ebe b938ca15 97b03cff 3dc2c0f8 8d1ab2ec 64380e51
68cc7bfb d90f2788 12490181 5de5ffd4 dd7ef86a 76a2e214 b9a40368 925d958f
4b39fffa ba39aee9 a4ffd30b faf7933b 6d498623 193cbcfa 27627545 825cf47a
61bd8ba0 d11e42d1 cead04f4 127ea392 10428db7 8272a972 9270c4a8 127de50b
285ba1c8 3c62f44f 35c0eaa5 e805d231 428929fb b4fcdf82 4fb66a53 0e7dc15b
1f081fab 108618ae fcfd086d f9ff2889 694bcc11 236a5cae 12deca4d 2c3f8cc5
d2d02dfe f8ef5896 e4cf52da 95155b67 494a488c b9b6a80c 5c8f82bc 89d36b45
3a609437 ec00c9a9 44715253 0a874b49 d773bc40 7c34671c 02717ef6 4feb5536
a2d02fff d2bf60c4 d43f03c0 50b4ef6d 07478cd1 006e1888 a2e53f55 b9e6d4bc
a2048016 97573833 d7207d67 de0f8f3d 72f87b33 abcc4f33 7688c55d 7b00a6b0
947b0001 570075d2 f9bb88f8 8942019e 4264a5ff 856302e0 72dbd92b ee971b69
6ea22fde 5f08ae2b af7a616d e5c98767 cf1febd2 61efc8c2 f1ac2571 cc8239c2
67214cb8 b1e583d1 b7dc3e62 7f10bdce f90a5c38 0ff0443d 606e6dc6 60543a49
5727c148 2be98a1d 8ab41738 20e1be24 af96da0f 68458425 99833be5 600d457d
282f9350 8334b362 d91d1120 2b6d8da0 642b1e31 9c305a00 52bce688 1b03588a
f7baefd5 4142ed9c a4315c11 83323ec5 dfef4636 a133c501 e9d3531c ee353783
```

S-Box $S_4$

```
9db30420 1fb6e9de a7be7bef d273a298 4a4f7bdb 64ad8c57 85510443 fa020ed1
7e287aff e60fb663 095f35a1 79ebf120 fd059d43 6497b7b1 f3641f63 241e4adf
28147f5f 4fa2b8cd c9430040 0cc32220 fdd30b30 c0a5374f 1d2d00d9 24147b15
ee4d111a 0fca5167 71ff904c 2d195ffe 1a05645f 0c13fefe 081b08ca 05170121
80530100 e83e5efe ac9af4f8 7fe72701 d2b8ee5f 06df4261 bb9e9b8a 7293ea25
ce84ffdf f5718801 3dd64b04 a26f263b 7ed48400 547eebe6 446d4ca0 6cf3d6f5
2649abdf aea0c7f5 36338cc1 503f7e93 d3772061 11b638e1 72500e03 f80eb2bb
abe0502e ec8d77de 57971e81 e14f6746 c9335400 6920318f 081dbb99 ffc304a5
4d351805 7f3d5ce3 a6c866c6 5d5bcca9 daec6fea 9f926f91 9f46222f 3991467d
a5bf6d8e 1143c44f 43958302 d0214eeb 022083b8 3fb6180c 18f8931e 281658e6
26486e3e 8bd78a70 7477e4c1 b506e07c f32d0a25 79098b02 e4eabb81 28123b23
69dead38 1574ca16 df871b62 211c40b7 a51a9ef9 0014377b 041e8ac8 09114003
```

```
bd59e4d2 e3d156d5 4fe876d5 2f91a340 557be8de 00eae4a7 0ce5c2ec 4db4bba6
e756bdff dd3369ac ec17b035 06572327 99afc8b0 56c8c391 6b65811c 5e146119
6e85cb75 be07c002 c2325577 893ff4ec 5bbfc92d d0ec3b25 b7801ab7 8d6d3b24
20c763ef c366a5fc 9c382880 0ace3205 aac9548a eca1d7c7 041afa32 1d16625a
6701902c 9b757a54 31d477f7 9126b031 36cc6fdb c70b8b46 d9e66a48 56e55a79
026a4ceb 52437eff 2f8f76b4 0df980a5 8674cde3 edda04eb 17a9be04 2c18f4df
b7747f9d ab2af7b4 efc34d20 2e096b7c 1741a254 e5b6a035 213d42f6 2c1c7c26
61c2f50f 6552daf9 d2c231f8 25130f69 d8167fa2 0418f2c8 001a96a6 0d1526ab
63315c21 5e0a72ec 49bafefd 187908d9 8d0dbd86 311170a7 3e9b640c cc3e10d7
d5cad3b6 0caec388 f73001e1 6c728aff 71eae2a1 1f9af36e cfcbd12f c1de8417
ac07be6b cb44a1d8 8b9b0f56 013988c3 b1c52fca b4be31cd d8782806 12a3a4e2
6f7de532 58fd7eb6 d01ee900 24adffc2 f4990fc5 9711aac5 001d7b95 82e5e7d2
109873f6 00613096 c32d9521 ada121ff 29908415 7fbb977f af9eb3db 29c9ed2a
5ce2a465 a730f32c d0aa3fe8 8a5cc091 d49e2ce7 0ce454a9 d60acd86 015f1919
77079103 dea03af6 78a8565e dee356df 21f05cbe 8b75e387 b3c50651 b8a5c3ef
d8eeb6d2 e523be77 c2154529 2f69efdf afe67afb f470c4b2 f3e0eb5b d6cc9876
39e4460c 1fda8538 1987832f ca007367 a99144f8 296b299e 492fc295 9266beab
b5676e69 9bd3ddda df7e052f db25701c 1b5e51ee f65324e6 6afce36c 0316cc04
8644213e b7dc59d0 7965291f ccd6fd43 41823979 932bcdf6 b657c34d 4edfd282
7ae5290c 3cb9536b 851e20fe 9833557e 13ecf0b0 d3ffb372 3f85c5c1 0aef7ed2
```

### *1.2 CAST-256 Notation*

The following notation is employed in the specification of CAST-256.

Let $f_1, f_2, f_3$ be as defined for CAST-128.

Let $\beta = (ABCD)$ be a 128-bit block where $A, B, C,$ and $D$ are each 32 bits in length.

Let "$\beta \leftarrow Q_i(\beta)$" be short-hand notation for the following:

$$
\begin{aligned}
C &= C \oplus f_1(D, k_{r_0}{}^{(i)}, k_{m_0}{}^{(i)}) \\
B &= B \oplus f_2(C, k_{r_1}{}^{(i)}, k_{m_1}{}^{(i)}) \\
A &= A \oplus f_3(B, k_{r_2}{}^{(i)}, k_{m_2}{}^{(i)}) \\
D &= D \oplus f_1(A, k_{r_3}{}^{(i)}, k_{m_3}{}^{(i)})
\end{aligned}
$$

Let "$\beta \leftarrow \overline{Q}_i(\beta)$" be short-hand notation for the following:

$$
\begin{aligned}
D &= D \oplus f_1(A, k_{r_3}{}^{(i)}, k_{m_3}{}^{(i)}) \\
A &= A \oplus f_3(B, k_{r_2}{}^{(i)}, k_{m_2}{}^{(i)}) \\
B &= B \oplus f_2(C, k_{r_1}{}^{(i)}, k_{m_1}{}^{(i)}) \\
C &= C \oplus f_1(D, k_{r_0}{}^{(i)}, k_{m_0}{}^{(i)})
\end{aligned}
$$

($Q(\cdot)$ is called a "forward quad-round" and $\overline{Q}(\cdot)$ is called a "reverse quad-round".)

Let $k_r{}^{(i)} = \{k_{r_0}{}^{(i)}, k_{r_1}{}^{(i)}, k_{r_2}{}^{(i)}, k_{r_3}{}^{(i)}\}$ be the set of rotation keys for the $i^{\text{th}}$ quad-round, where $k_{r_j}{}^{(i)}$ is a 5-bit rotation key for $f_1, f_2,$ or $f_3$ (as specified above).

Let $k_m{}^{(i)} = \{k_{m_0}{}^{(i)}, k_{m_1}{}^{(i)}, k_{m_2}{}^{(i)}, k_{m_3}{}^{(i)}\}$ be the set of masking keys for the $i^{\text{th}}$ quad-round, where $k_{m_j}{}^{(i)}$ is a 32-bit masking key for $f_1, f_2,$ or $f_3$ (as specified above).

*CAST-256 Notation (cont'd)*

Let $\kappa = (ABCDEFGH)$ be a 256-bit block where $A, B, \ldots, H$ are each 32 bits in length.

Let "$\kappa \leftarrow \omega_i(\kappa)$" be short-hand notation for the following:

$$
\begin{aligned}
G &= G \oplus f_1(H, t_{r_0}{}^{(i)}, t_{m_0}{}^{(i)}) \\
F &= F \oplus f_2(G, t_{r_1}{}^{(i)}, t_{m_1}{}^{(i)}) \\
E &= E \oplus f_3(F, t_{r_2}{}^{(i)}, t_{m_2}{}^{(i)}) \\
D &= D \oplus f_1(E, t_{r_3}{}^{(i)}, t_{m_3}{}^{(i)}) \\
C &= C \oplus f_2(D, t_{r_4}{}^{(i)}, t_{m_4}{}^{(i)}) \\
B &= B \oplus f_3(C, t_{r_5}{}^{(i)}, t_{m_5}{}^{(i)}) \\
A &= A \oplus f_1(B, t_{r_6}{}^{(i)}, t_{m_6}{}^{(i)}) \\
H &= H \oplus f_2(A, t_{r_7}{}^{(i)}, t_{m_7}{}^{(i)})
\end{aligned}
$$

($\omega(\cdot)$ is called a "forward octave".)

Let "$k_r{}^{(i)} \leftarrow \kappa$" be short-hand notation for the following:

$$k_{r_0}{}^{(i)} = 5LSB(A),\ \ k_{r_1}{}^{(i)} = 5LSB(C),\ \ k_{r_2}{}^{(i)} = 5LSB(E),\ \ k_{r_3}{}^{(i)} = 5LSB(G)$$

where $5LSB(x)$ denotes "the five least significant bits of *x*".

Let "$k_m{}^{(i)} \leftarrow \kappa$" be short-hand notation for the following:

$$k_{m_0}{}^{(i)} = H,\ \ k_{m_1}{}^{(i)} = F,\ \ k_{m_2}{}^{(i)} = D,\ \ k_{m_3}{}^{(i)} = B$$

### *1.3 The CAST-256 Cipher*

$\beta =$ 128 bits of plaintext.

```
$for(i = 0;\ i < 6;\ i++)$
        $\beta \leftarrow Q_i(\beta)$
$for(i = 6;\ i < 12;\ i++)$
        $\beta \leftarrow \overline{Q}_i(\beta)$
```

128 bits of ciphertext $= \beta$

### *Round Key Re-Ordering for Decryption*

The cipher employs a 256-bit primary key $K$. Decryption is identical to encryption except that the sets of quad-round keys $k_r^{(i)}, k_m^{(i)}$ derived from $K$ are used in reverse order as follows.

```
$for(i = 0; i < 12; i++)\{$
        $k_{r_{new}}^{(i)} = k_r^{(11-i)}$
        $k_{m_{new}}^{(i)} = k_m^{(11-i)}$
$\}$
```

## *1.4 The CAST-256 Key Schedule*

*Initialization:*

$$c_m = 2^{30}\sqrt{2} = 5A827999_{16}$$
$$m_m = 2^{30}\sqrt{3} = 6ED9EBA1_{16}$$
$$c_r = 19$$
$$m_r = 17$$

```
$for(i = 0;\ i < 24;\ i++)$
    $for(j = 0;\ j < 8;\ j++)\{$
        $t_{m_j}^{(i)} = c_m$
        $c_m = (c_m + m_m) \bmod 2^{32}$
        $t_{r_j}^{(i)} = c_r$
        $c_r = (c_r + m_r) \bmod 32$
    $\}$
```

*Key Schedule:*

$\kappa = ABCDEFGH =$ 256 bits of primary key, $K$.

```
$for(i = 0;\ i < 12;\ i++)\{$
    $\kappa \leftarrow \omega_{2i}(\kappa)$
    $\kappa \leftarrow \omega_{2i+1}(\kappa)$
    $k_r^{(i)} \leftarrow \kappa$
    $k_m^{(i)} \leftarrow \kappa$
$\}$
```

Note:

$$(|K| = 128) \Rightarrow (E = F = G = H = 0)$$
$$(|K| = 160) \Rightarrow (F = G = H = 0)$$
$$(|K| = 192) \Rightarrow (G = H = 0)$$
$$(|K| = 224) \Rightarrow (H = 0)$$

## 2. Design Rationale

### *2.1 Overall Structure*

The fundamental mechanism for the expansion of a 64-bit block size to a larger block size is the generalization of the basic Feistel network (Schneier and Kelsey [SK96] have referred to the structure used here as an "incomplete" Feistel network). The motivation is as follows. In a traditional Feistel network (such as DES), rather than thinking of the exchange of left and right halves in each round as a "swap", it may be viewed as a circular right-shift of 32 bits. Such a view allows one to consider a cipher with a block size of 32*n* bits, which uses the same round function as the original cipher but requires *n* rounds (instead of 2) to input all bits of the block to the round function.

A picture may help to clarify the operation.

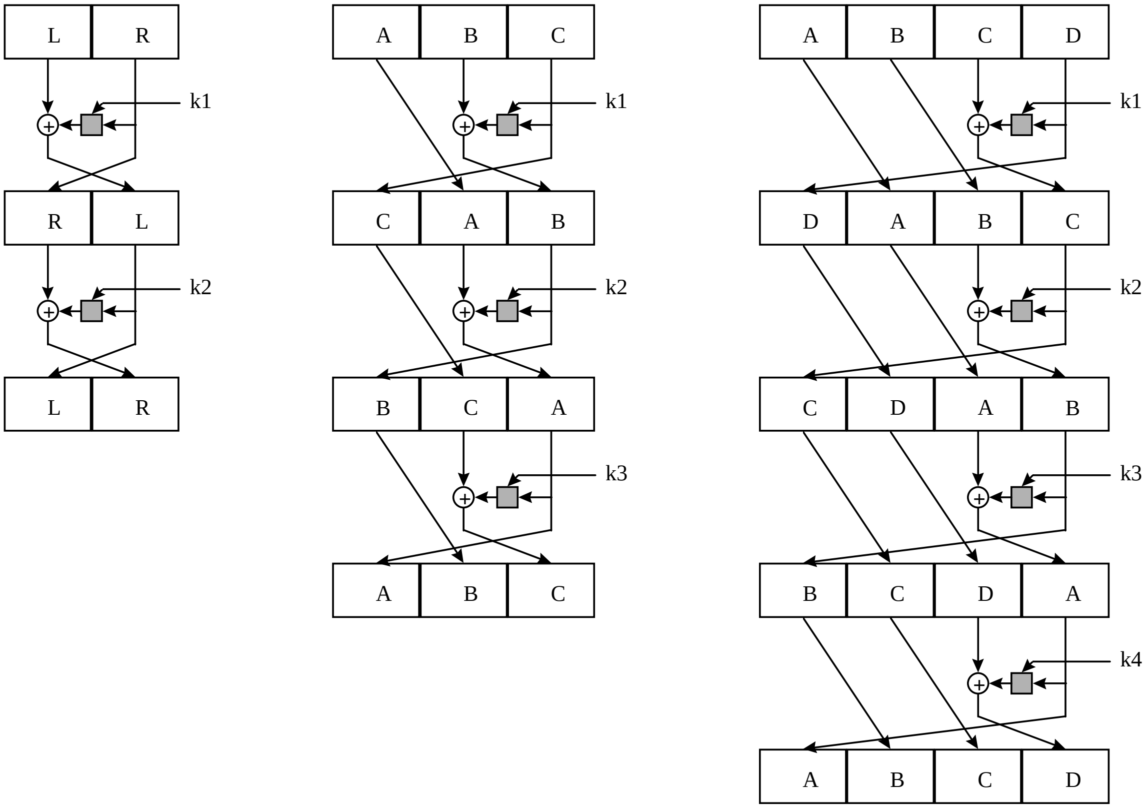

***Figure 1***

The left-most diagram is the “traditional” Feistel network. If this describes two rounds of DES, then *L* and *R* are each 32 bits in length and the cipher has a 64-bit block size. Continuing the illustration, the middle diagram describes an extended Feistel network for a cipher with a 96-bit block size, and the right-most diagram describes the structure of a cipher with a 128-bit block size. In each case, we may think of the number of rounds shown as a basic “unit” (in terms of submitting all input bits to the round function); the actual number of rounds chosen for the full cipher will be some multiple of this “unit” (e.g., for DES, the multiple is 8).

### *2.2 Decryption Considerations*

The disadvantage of the generalized structure given above is that it requires a separate structure for decryption (since data must be left-shifted, rather than right-shifted, in each round in order to go backwards through the rounds). By contrast, with the “traditional” Feistel network decryption and encryption are identical except for a change in the ordering of the round keys so no separate structure is needed. Clearly, in constrained environments (such as hardware or firmware implementations that are very resource-limited) requiring two structures is unattractive.

A simple solution to the above concern is to design the structure such that if there are *r* rounds in the full cipher, the first *r*/2 rounds use right-shifting (as shown in the diagram above) and the last *r*/2 rounds use left-shifting. In this way, the desirable feature of “traditional” Feistel networks with respect to decryption (i.e., that decryption is identical to encryption, requiring only a reversal of the round keys) is preserved. This simplifies implementation and operation of the cipher and helps to make its use feasible in resource-limited environments.

### *2.3 Choice of Round Function*

One of the very attractive features of the generalized structure given above is that it enables direct re-use of the round function from the “traditional” Feistel network. Within the class of DES-like ciphers, it is well known that increasing the size of the round function typically involves increasing the size of its component substitution boxes (s-boxes); it is also well known that increasing s-box size is generally difficult. For those ciphers that already employ large s-boxes, size increases can be a monumental task. [As a particular example, doubling the input and output sizes of a carefully-constructed 8×32 s-box may require a work factor of roughly $2^{64}$ steps (more than is necessary to break DES by exhaustive search!), aside from the fact that the resulting s-box grows from 4 Kbytes to more than half a million bytes of memory.] Being able to re-use the original round function is therefore very desirable. The important technical decision, however, is which “traditional” Feistel network round function to use in the generalized network.

The CAST-128 set of round functions has a number of appealing features.

- Firstly, the component bent-function-based s-boxes are designed according to a mathematical procedure which produces substitution boxes with several important cryptographic properties (such as high nonlinearity, low XOR difference distribution table values, good higher-order Strict Avalanche Criterion, and good higher-order (Output) Bit Independence Criterion) [A97b].

- Secondly, the use of both a "masking" key and a "rotation" key ensures that the key entropy is higher than the data entropy in each round (following the recommendation of [RPD97]) and appears to make the construction of iterative statistical attacks such as linear and differential cryptanalysis significantly more difficult (or impossible) [A97b].

- Thirdly, the mixing of operations from different algebraic groups (addition modulo 2 and addition / subtraction modulo $2^{32}$) appears to be effective not only in reducing the probability of the round differential [AM97, O'C98], but in reducing the possibility of higher-order differential attacks as well [MSK98].

- Finally, mixing the order of the group operations (i.e., by varying the order of round functions throughout the cipher, as is done in CAST-128) appears to frustrate the practical construction of iterative characteristics.

In summary, then, the extensive analysis done on the CAST design procedure (including focused attention within several master's- and doctoral-level theses on symmetric cipher design and analysis) lends confidence to its choice as the round function for this generalized Feistel network.

[See *CAST-256: Algorithm Analysis* below for a partial list of published work which discusses or analyzes various aspects of the CAST design procedure. For one significant example of unpublished work that has been done on CAST, the Communications Security Establishment, after extensive analysis, has determined and will formally state that the CAST-128 algorithm is suitable for the protection of all levels of Designated information within the Government of Canada. Please see the attached letter dated June 5th, 1998, and note that "CAST5" is the name used for "CAST-128" when specific key lengths are explicitly intended (see [A97c], Section 2.5).

### *2.4 Number of Rounds*

Given that the basic unit (see “Overall Structure” above) in DES is a “double round” and that a multiple of 8 is used to give the full 16-round cipher, it is reasonable to conclude that a 128-bit block size, with a “quad-round” as the basic unit, should consist of at least 32 rounds for the full cipher. It is important to note, however, that a cipher being constructed as a candidate for AES consideration must support not only twice the block size of CAST-128, but twice the key size as well. A deeper security analysis (see attached document, *CAST-256: Algorithm Analysis*) suggests that 48 rounds (i.e., 12 “quad rounds”) provides security protection commensurate with the parameters of the desired cipher.

### *2.5 Key Schedule*

Key scheduling (deriving a set of round keys from an initial key) is an extremely important aspect of cipher design since sub-optimal key schedules can lead to exploitable weaknesses in the cipher (including weak keys, equivalent keys, complementation properties, and susceptibility to related-key attacks), and overly-complicated key schedules can lead to prohibitively-long set-up times (limiting the use of the cipher in some environments).

The design philosophy chosen for the CAST-256 key schedule is identical to that chosen for the CAST-256 cipher itself: the key schedule essentially describes a generalized Feistel network with a 256-bit block size. A simple (but fixed) set of round keys is used to key this network and the CAST-256 initial key is used as the plaintext input. Some of the output bits of selected rounds during this “encryption” define the actual round keys for the CAST-256 cipher. Important features of this key scheduling approach include the following.

- The inherent strength of the generalized Feistel network is used in the key schedule to create round keys, increasing confidence that the set of key values (comprised of the generated round keys and the CAST-256 initial key) will appear to be pair-wise independent to any statistical analysis.

- If an attack can be mounted that derives four or more full round keys (i.e., full masking keys and the corresponding rotation keys) from the CAST-256 cipher, it still appears to require a computational effort of at least $2^{256 - (4 * 32) - (4 * 5)} = 2^{108}$ guesses to derive the CAST-256 initial key from this information.

- Since the key schedule describes a generalized Feistel network, it is extremely unlikely that key collisions can occur. The key schedule defines a cipher with a fixed key (i.e., a permutation over the input space) so for two different CAST-256 initial keys to produce identical sets of round keys, the different cipher inputs would have to map to round function outputs (in every relevant round) that differed only in the 108 bits *not* used to produce round key bits. The probability of this occurring in each octave that produces round keys is $2^{108}/2^{256} = 2^{-148}$, so the probability that this occurs over the full set of round keys is $2^{-148*12} = 2^{-1776}$ (essentially zero, since there are only $2^{256}$ possible initial keys).

- The key scheduling operation requires the equivalent of four CAST-256 encryption operations to produce a full set of round keys. This ratio is not prohibitive for most environments and compares favorably with many current implementations of DES.

The key schedule chosen for CAST-256 appears to have a number of desirable cryptographic features and takes into account much of the research into key schedule design and analysis over the past two decades (see, for example, [A94] and the references included in [A97]).

### *2.6 Conclusions*

A number of alternatives exist for doubling the block size of a cipher from 64 bits to 128 bits, including the following.

- Feistel network. In such a design, the round function of the Feistel network is the original 64-bit cipher, which may itself be a Feistel network (this is a simple extension of ideas presented in, for example, [LR88, L96]).

- Substitution-Permutation (SP) network [F73]. In such a design, two parallel implementations of the original cipher are used as the substitution layers; these are interspersed with an extended permutation layer (i.e., a permutation which is the width of the desired block size).

- "Fenced" Construction [R96]. In such a design, two parallel implementations of the original cipher are surrounded by specially-constructed mixing operations, which in turn are surrounded by a layer of substitution boxes.

However, it was felt that all the alternatives considered had one or more drawbacks which made them somewhat less attractive as AES submission candidates. For example, the Feistel network suffers significant security degradation if one or two rounds may be "peeled off" by some attack (not an uncommon situation) since the entire outer network would likely consist of only four or six rounds (for performance reasons). The SP network may be subject to poor encryption / decryption performance since even two substitution layers with a permutation layer in between (the minimum possible configuration) halves the speed of the original cipher; a larger number of layers decreases performance significantly beyond this. Finally, the Fenced construction has non-trivial design and implementation impacts with the need for solid theoretical justification for the particular mixing operations used and the need for sufficient processing time and memory for the pseudo-random generation and storage of the necessary s-boxes.

The approach taken in this proposal to achieve block size doubling (i.e., the use of a generalized Feistel network) appears to be the simplest and most elegant of the various alternatives. It has none of the drawbacks listed above, is straightforward to understand and to analyze, and builds on the confidence gained from the extensive literature on ciphers based on Feistel networks. Furthermore, it allows unmodified re-use of a round function with a number of attractive cryptographic features, and suggests an intuitive architecture for the associated key scheduling algorithm.

We conclude that the rationale for CAST-256 is solid, resting on firm theoretical results and immediately appealing, defensible, concepts for every aspect of the cipher design. The resulting algorithm has good performance, reasonable code and memory size, and high security (according to all analysis conducted to date); it thus appears to meet all the requirements for an AES submission candidate.

## 3. Bit Naming / Numbering Convention Provided

True (needed only in section *1.1 CAST-128 Notation* above, where most- to least-significant bytes of a 32-bit word are specified).

## 4. No Parity Bits Specified in the Key Definition

True.

## 5. References

[A94] C. Adams, “Simple and Effective Key Scheduling for Symmetric Ciphers”, in *Workshop Record* of the Workshop on Selected Areas in Cryptography (SAC ’94), Kingston, Canada, May 1994, pp.129-133.

[A97] C. Adams, “DES-80”, in *Workshop Record* of the Workshop on Selected Areas in Cryptography (SAC ’97), Ottawa, Canada, August 1997, pp.160-171.

[A97b] C. Adams, “Constructing Symmetric Ciphers Using the CAST Design Procedure”, in *Selected Areas in Cryptography*, E. Kranakis and P. Van Oorschot (ed.), Kluwer Academic Publishers, 1997, pp.71-104 (reprinted from *Designs, Codes and Cryptography*, vol. 12, no. 3, November 1997).

[A97c] C. Adams, “The CAST-128 Encryption Algorithm”, RFC 2144, May 1997.

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[L96] S. Lucks, “Faster Luby-Rackoff Ciphers”, in *Proceedings of the Third International Workshop on Fast Software Encryption*, Cambridge, UK, February 1996, Springer, LNCS 1039, pp.189-203.

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[O’C98] L. O’Connor, Preliminary analytical results concerning the mixing of operations from different algebraic groups and the maximum value of the resulting XOR difference distribution table (unpublished).

[R96] T. Ritter, “The Fenced DES Cipher: Stronger Than DES But Made From DES”, *http://www.io.com/~ritter/FENCED.HTM*, November 10, 1996.

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[SK96] B. Schneier and J. Kelsey, “Unbalanced Feistel Networks and Block Cipher Design”, in *Proceedings of the Third International Workshop on Fast Software Encryption*, Cambridge, UK, February 1996, Springer, LNCS 1039, pp.121-144.

Communications Security
Establishment

P.O. Box 9703
Terminal
Ottawa, Canada
K1G 3Z4

Centre de la sécurité
des télécommunications

C.P. 9703
Terminus
Ottawa, Canada
K1G 3Z4

05 June 1998

Mr. Brian O'Higgins
Executive Vice President and
Chief Technology Officer
Entrust Technologies
750 Heron Road
Suite 800
Ottawa, Ontario
K1V 1A7

Dear Mr. O'Higgins,

I am very pleased to advise you that CSE has completed its evaluation of the CAST5 algorithm (80 and 128 bit versions). We have determined that CAST5 is suitable for the protection of all levels of Designated information within the GOC. A formal statement of this approval will be promulgated to Government of Canada departments and agencies in the very near future.

On behalf of the Communications Security Establishment please accept my congratulations.

David McKerrow
Communications Security Establishment
Director General Information Technology Security

# CAST-256
# Computational Efficiency

## 1. Efficiency Estimates for the NIST AES Analysis Platform

### *1.1 Platform Description*

IBM-compatible PC, with an Intel Pentium Pro Processor, 200MHz clock speed, 64MB RAM, running Windows95.

### *1.2 Speed Estimates (in clock cycles)*

| Operation | 128/128 | 192/128 | 256/128 |
|---|---|---|---|
| Encrypt one data block: | 1790 | 1790 | 1790 |
| Decrypt one data block: | 1790 | 1790 | 1790 |
| Key setup: | 9090 | 9090 | 9090 |
| Algorithm setup: | 0 | 0 | 0 |
| Key change: | 9090 | 9090 | 9090 |

### *1.3 Tradeoffs Between Speed and Memory*

For environments in which memory is not a scarce resource, s-boxes $S_1$ and $S_2$ can be combined into three 16×32 s-boxes (one corresponding to $S_1 \oplus S_2$, one corresponding to $S_1$-$S_2$, and one corresponding to $S_1$+$S_2$, for each of the three round function types). This saves one memory lookup and combining operation per round, which will result in a modest performance increase.

## 2. Efficiency Estimates for 8-Bit Processors

### *2.1 Platform Description*

Motorola 6811 microprocessor, 2MHz clock speed, assembly language implementation.

### *2.2 Speed Estimates (in clock cycles)*

| Operation | 128/128 | 192/128 | 256/128 |
|---|---|---|---|
| Encrypt one data block: | 26000 | 26000 | 26000 |
| Decrypt one data block: | 26000 | 26000 | 26000 |
| Key setup: | 110000 | 110000 | 110000 |
| Algorithm setup: | 0 ms | 0 ms | 0 ms |
| Key change: | 110000 | 110000 | 110000 |

### *2.3 Tradeoffs Between Speed and Memory*

None known.

## 3. Efficiency Estimates for Other Platforms

### *3.1 Platform Description*

IBM-compatible PC, with an Intel Pentium II Processor, 300MHz clock speed, 128MB RAM, running Windows NT 4.0, assembly language implementation.

### *3.2 Speed Estimates (in clock cycles)*

| Operation | 128/128 | 192/128 | 256/128 |
|---|---|---|---|
| Encrypt one data block: | 815 | 815 | 815 |
| Decrypt one data block: | 815 | 815 | 815 |
| Key setup: | 4130 | 4130 | 4130 |
| Algorithm setup: | 0 | 0 | 0 |
| Key change: | 4130 | 4130 | 4130 |

### *3.3 Tradeoffs Between Speed and Memory*

For environments in which memory is not a scarce resource, s-boxes $S_1$ and $S_2$ can be combined into three 16×32 s-boxes (one corresponding to $S_1 \oplus S_2$, one corresponding to $S_1$-$S_2$, and one corresponding to $S_1$+$S_2$, for each of the three round function types). This saves one memory lookup and combining operation per round, which will result in a modest performance increase.

## 4. Efficiency Estimates for Other Platforms

### *4.1 Platform Description*

Sun UltraSparc 1, 167MHz clock speed, 124MB RAM, running Solaris 2.5.

### *4.2 Speed Estimates (in clock cycles)*

| Operation | 128/128 | 192/128 | 256/128 |
|---|---|---|---|
| Encrypt one data block: | 1180 | 1180 | 1180 |
| Decrypt one data block: | 1180 | 1180 | 1180 |
| Key setup: | 5830 | 5830 | 5830 |
| Algorithm setup: | 0 | 0 | 0 |
| Key change: | 5830 | 5830 | 5830 |

### *4.3 Tradeoffs Between Speed and Memory*

For environments in which memory is not a scarce resource, s-boxes $S_1$ and $S_2$ can be combined into three 16×32 s-boxes (one corresponding to $S_1 \oplus S_2$, one corresponding to $S_1 - S_2$, and one corresponding to $S_1 + S_2$, for each of the three round function types). This saves one memory lookup and combining operation per round, which will result in a modest performance increase.

## 5. General Efficiency Comments

As will be noted in the tables given above, CAST-256 has the following features:

- it requires no algorithm setup time (e.g., there is no need to generate s-boxes or other tables, and no need to pre-compute values);
- decryption performance is identical to encryption performance;
- key change time is identical to key setup time;
- there is no penalty for key size differences (i.e., encryption / decryption performance and key setup performance are unaffected by whether the primary key is 128 bits, 256 bits, or a value in between).

# CAST-256
# Algorithm Analysis

## 1. Analysis With Respect to Known Attacks

The classical attacks on ciphers are as follows: *ciphertext only*; *known plaintext*; and *chosen plaintext*. The advent of public-key cryptography added utility to the concept of a *chosen ciphertext* attack, but this appears to be of little added value in the analysis of symmetric ciphers. Research in the past decade or so has also introduced the notions of *chosen key* and *related key* attacks, which have enjoyed some success in the cryptanalysis of specific symmetric ciphers. Within the iterated symmetric ciphers (the class of algorithms to which CAST-256 belongs), the techniques known as *linear cryptanalysis* and *differential cryptanalysis* (along with their combinations and higher-orders) currently represent the most general and powerful instances of *known plaintext* and *chosen plaintext* attacks, respectively.

This section of the submission package examines the CAST-256 algorithm with respect to the families of cryptanalytic attack listed above.

### *1.1 Ciphertext Only Attack*

No techniques are currently known that will allow an attacker to infer or derive information about the plaintext, the primary key, or any subset of round keys from any collection of ciphertext blocks. The one (unavoidable) exception to this is the technique applicable to all $n$-bit-block ciphers when used in Cipher-Block-Chaining (CBC) mode: once $2^{n/2}$ blocks have been encrypted, with probability roughly ½ (rapidly increasing as more blocks are encrypted) an XOR relationship between a particular pair of plaintexts will be known.

### *1.2 Known Plaintext Attack: Linear Cryptanalysis*

Linear cryptanalysis [M94] attempts to exploit any high-probability occurrences of linear expressions of input, output, and round key bits in the round function of an iterated cipher. It has been approximated [M94] that the best linear expression for $r$-rounds of a cipher has a probability of being satisfied that is bounded as follows:

$$\left|p_L - \tfrac{1}{2}\right| \le 2^{\alpha-1} \cdot \left|p_\beta - \tfrac{1}{2}\right|^{\alpha}$$

where $p_L$ represents the probability that the linear expression holds, $p_\beta$ represents the probability of the best linear approximation, and $\alpha$ represents the number of s-boxes involved in that linear approximation. This expression is based on the assumption of independent round keys such that the linear approximations of the s-boxes are independent. In an analogous way to “differentials” and “characteristics” in differential cryptanalysis, provable immunity in linear cryptanalysis relies on bounding the likelihood of an overall linear expression (sometimes referred to as the “linear hull”) rather than any particular linear “characteristic”. However, for many ciphers (including CAST-256) this is a difficult analytical task. What are typically considered, therefore, are the building blocks of an overall linear expression: the sequence of approximations of the round functions which result in the overall linear expression.

A basic linear attack typically uses a sequence of linear approximations of the rounds to create an overall linear expression involving subsets of plaintext and ciphertext bits. From this it is possible to derive the equivalent of one key bit represented as the XOR of a number of round key bits. In this case, it is shown [M94] that the number of known plaintexts required is approximately

$$N_L = \left|p_L - \tfrac{1}{2}\right|^{-2} .$$

It can be shown that the best linear approximation has a probability given by

$$\left|p_\beta - \tfrac{1}{2}\right| = \frac{2^{m-1} - NL_{\min}}{2^m}$$

where $m$ is the number of input bits to the s-box and $NL_{\min}$ is the nonlinearity of the s-box [LHT97]. For the s-boxes of CAST-256, $m = 8$ and $NL_{\min} = 74$. Furthermore, for the CAST-256 cipher, the best linear approximation appears to involve 4 s-boxes every 4 rounds such that the linear approximation of the round function for every 4th round involves no output bits. That is, the linear expression used is $X_{i_1} \oplus X_{i_2} \oplus \ldots \oplus X_{i_t}$, where $X_{i_j}$ represents an input bit to the s-box. Hence, for an $r$-round linear approximation, $\alpha = r$. The number of known plaintexts required for a 48-round linear approximation of CAST-256, then, is approximately $2^{122}$. Note that this is almost equal to the total number of plaintexts available ($2^{128}$) and argues against the practicality of a linear attack on this cipher.

Furthermore, Youssef, *et al*, have proposed [YCT97] that a more accurate bound on the number of plaintexts required for linear cryptanalysis of a CAST cipher can be obtained by considering the combination of s-boxes in the round function rather than the individual s-boxes. In particular, they compute the value for $NL_S$, the nonlinearity of the composite

32×32 s-box when the individual 8×32 s-boxes are combined using XOR. Using this in place of $NL_{\min}$ in the equations above and setting $m = 32$ and $\alpha = r/2$ (since an $r$-round linear approximation must involve at least as many 32×32 s-boxes as $r/2$ iterations of the best 2-round approximation) yields a number of known plaintexts required for a 48-round linear approximation at more than $2^{174}$ (far beyond the number of plaintexts available). Note that experimental evidence suggests that combining s-boxes using mixed operations may increase the nonlinearity of the composite s-box even further.

It therefore appears that CAST-256 is immune to a linear cryptanalysis attack.

### *1.3 Chosen Plaintext Attack: Differential Cryptanalysis*

Differential cryptanalysis [BS93] attempts to exploit any high-probability output differences resulting from particular input differences in the round function of an iterated cipher. A block cipher can be proved to be resistant to differential cryptanalysis if it can be shown that no high-probability differentials [LMM91] exist, where an *i*-round differential is defined to be the XOR of two outputs after *i* rounds, where the outputs correspond to two plaintexts with a given XOR.

In a good cipher the probability of all differentials should approach $2^{-N}$, where $N$ is the block size. Strictly speaking, differential cryptanalysis requires only the existence of a highly-probable differential to succeed. However, differentials can be viewed to be comprised of a number of possible characteristics, where a characteristic specifies the exact sequence of input and output XORs for each round to achieve the overall differential input and output XOR.

It is typically difficult to derive the probability of any particular differential and, in practice, it would be hard for a cryptanalyst to determine the existence of a highly-probable differential without searching for highly-probable characteristics. Although it is often the case that an upper bound on the probability of a differential cannot be stated for a particular cipher (that is, resistance to a differential cryptanalytic attack cannot be proved), the probabilities of the most likely characteristics can be determined. These probabilities can then be used as a measure of the cipher’s resistance to differential cryptanalysis.

As is common in the literature, the analysis here is based on the assumption that all round keys are independent (although this assumption is not always necessary; see [C97]) and that the occurrence of output XORs given particular input XORs is independent for different rounds. Under such conditions, the probability of an $r$-round characteristic is given by

$$p_{\Omega_r} = \prod_{i=1}^{r} p_i$$

where $p_i$ represents the probability of the output XOR given the input XOR in round *i*. The best characteristics that can be constructed are typically iterative in nature. For the CAST-256 cipher with *R* rounds, the following appears to be the best possible *r*-round characteristic, where *r* is a multiple of 4. (Note that the notation (*W*,*X*,*Y*,*Z*) represents XOR vectors for the four 32-bit sub-blocks in a CAST-256 round function input.)

| | |
|---|---|
| $(0,0,0,\Delta)$ | [input XOR to round 1] |
| $0 \leftarrow \Delta$ with probability $p$ | [round 1] |
| $0 \leftarrow 0$ with probability 1 | [round 2] |
| $0 \leftarrow 0$ with probability 1 | [round 3] |
| $0 \leftarrow 0$ with probability 1 | [round 4] |
| ... | repeat up to $R/2$ rounds |
| $(0,\Delta,0,0)$, or some variation | [input XOR to round $(R/2 + 1)$] |
| $0 \leftarrow 0$ with probability 1 | [round $(R/2 + 1)$] |
| $0 \leftarrow 0$ with probability 1 | [round $(R/2 + 2)$] |
| $0 \leftarrow \Delta$ with probability $p$ | [round $(R/2 + 3)$] |
| $0 \leftarrow 0$ with probability 1 | [round $(R/2 + 4)$] |
| ... | repeat up to *r* rounds for *r*-round char. |

The input XOR to round ($R/2$ +1) will be a vector in which one of the sub-blocks is non-zero and the other three sub-blocks are zero (the precise variation which applies for a given cipher depends upon the value of *R*). Without loss of generality, the example $(0,\Delta,0,0)$ is shown above.

As per the analysis and rationale given in [LHT97], the input-output XOR pair for a simplified CAST round function (i.e., one which does not include the key-dependent rotation, and for which the only s-box combining operation used is XOR) can be assumed to have a probability of $p \le 2^{-14}$. This is based on the fact that all four s-boxes in the CAST round function are injective and the format of the XOR pair has the output XOR being equal to 0. This leads to the conclusion that the best *r*-round iterated characteristic as shown above has a probability given by

$$p_{\Omega_r} \le (2^{-14})^{r/4}$$

In particular, a 40-round characteristic must have a probability less than or equal to $2^{-140}$ according to the assumptions of the analysis. This implies that the number of chosen plaintexts required for this attack would be greater than $2^{140}$ for the 48-round cipher (substantially greater than the number of plaintexts available for a 128-bit block size).

It therefore appears that CAST-256 is immune to a differential cryptanalysis attack.

### *1.4 Chosen Key Attack*

CAST-256 appears to be secure with respect to this attack. The use of a cipher (built around the CAST-128 set of round functions) as a key schedule gives confidence that no exploitable statistical correlation exists between the primary key and the set of generated round keys. Thus, allowing an attacker to choose a particular primary key difference appears to yield no exploitable similarities in the corresponding sets of round keys compared with the victim encrypting with two randomly-chosen primary keys.

### *1.5 Related Key Attack*

CAST-256 appears to be secure with respect to this attack. The use of a cipher (built around the CAST-128 set of round functions) as a key schedule gives confidence that no exploitable statistical correlations exist within the set of generated round keys. Thus, this attack, which depends upon the use of a simple derivation algorithm for a round key from previous round keys, appears not to be applicable to CAST-256.

### *1.6 Enhancements to the Above Statistical Attacks: Combinations and Higher-Orders*

The analysis given above for both linear and differential cryptanalysis applies to a greatly simplified version of the CAST-256 cipher. The actual cipher, which includes key-dependent rotation and mixed operations in the round function (both for data masking and for s-box combination), appears to be much more difficult / impossible to attack using the methods as described in [M94] and [BS93] (see [A97] for some discussion of this). In particular, experiments in which two CAST-256 s-boxes are combined using addition or subtraction modulo $2^{32}$ show that the maximum value in the XOR difference distribution table is approximately 10% of the maximum that occurs when the s-boxes are combined using XOR. Experiments on combinations of three CAST-256 s-boxes are on-going, but thus far show similar results. This lends confidence that combinations of four s-boxes using mixed operations (as is done in the CAST-256 round function) are effective in increasing resistance to differential cryptanalysis.

The above experimental work [AM97] is supported by a new analytical result [O'C98], which shows that for a random $n$-bit permutation, the probability that the maximum entry in a differential table based on XOR differences is greater than a bound $B_n$ approaches 1 as $n$ grows, whereas the probability that the maximum entry in a table based on non-XOR differences (e.g., modular addition or multiplication) is greater than that same bound approaches 0. Furthermore, the bound is accurate for the 8-bit case. Thus, although the details of the analyzed structure differ slightly from the internals of the CAST-256 round function as used in the above experiments, the conclusion is the same: using operations from different algebraic groups appears to be helpful in increasing resistance to differential cryptanalysis (by lowering the differential probability of a single round).

### 1.6.1 Combination Attacks

CAST-256 appears to be immune to both linear and differential cryptanalysis (requiring more plaintext than is available from the 128-bit block size) and appears to be immune to both chosen and related key attacks (due to the absence of exploitable statistical correlations among its generated keys). Given this, it seems highly unlikely that various combination attacks (such as *linear-differential*, or *differential-related-key*) can have any measure of success.

It therefore appears that this cipher is immune to the combination attacks currently known in the literature.

### 1.6.2 Higher-Order Attacks

The concept of *higher-order differentials* has been introduced [L94, K95] and used to successfully cryptanalyze ciphers proved secure against ordinary differential cryptanalysis [JK97]. A simplified version of the CAST-128 cipher (one which uses XOR for all operations in the round function) has been examined with respect to the higher-order differential attack [MSK98]. It has been shown that this attack is successful up to 5 rounds, but cannot be extended to higher numbers of rounds. Furthermore, the introduction of the key-dependent rotation operation is effective in increasing the computational complexity of this attack. Finally, the use of operations from different algebraic groups “makes the degree too high to cryptanalyze by the higher-order differential attack” [MSK98], so that the attack cannot even be mounted on a 5-round version of the cipher.

It therefore appears that CAST-256 (which has 48 rounds and uses the CAST-128 round functions) is immune to a higher-order differential attack.

## 2. Statements Regarding Properties of Keys

This section provides statements regarding the following properties of keys with respect to CAST-256: *weak keys*, *semi-weak keys*, *fixed points of a key*, *equivalent keys*, and *restrictions on key selection*. It also includes a statement on *complementation properties* since this is sometimes related to the way that round keys are used within a DES-like cipher.

### 2.1 Weak Keys

None known. In the CAST-256 cipher, all keys appear to be of equivalent strength and are usable for double encryption (i.e., no key appears to be its own inverse).

### 2.2 Semi-Weak Keys

None known. In the CAST-256 cipher, there appear to be no pairs of keys which cannot be used for double encryption (i.e., there do not appear to be pairs of keys $k_i$ and $k_j$ such that $k_j$ is the inverse of $k_i$).

### 2.3 Fixed Points of a key K

None known. From all evidence available thus far in the open literature, fixed points have only been easily found (i.e., requiring a level of effort for an $n$-bit block cipher of roughly $2^{n/2}$ operations rather than $2^n$ operations) in DES-like ciphers for weak and semi-weak keys. It therefore appears that CAST-256 has no easily-found fixed points for any key.

### 2.4 Equivalent Keys

None known. The key schedule defines a cipher with a fixed key (i.e., a permutation over the input space) so for two different CAST-256 initial keys to produce identical sets of round keys, the different cipher inputs would have to map to round function outputs (in every relevant round) that differed only in the 108 bits *not* used to produce round key bits. The probability of this occurring in each octave that produces round keys is $2^{108}/2^{256} = 2^{-148}$, so the probability that this occurs over the full set of round keys is $2^{-148*12} = 2^{-1776}$ (essentially zero, since there are only $2^{256}$ possible initial keys).

### *2.5 Restrictions on Key Selection*

None known. The key scheduling algorithm defines a symmetric block cipher with a fixed key where the CAST-256 primary key is used as the plaintext input. Because in this symmetric block cipher there are no restrictions on the input space (i.e., any plaintext can be encrypted), it follows that no restrictions are placed upon selection of CAST-256 primary keys.

### *2.6 Complementation Properties*

None known. There appear to be no complementations of combinations of plaintext, key, and ciphertext that lead to identities. This is due to the complexity of the key scheduling operation (so that complementing the primary key leads to random-looking changes to all round keys) and also to the use of multiple operations to combine data, key, and s-boxes within the round functions (XOR, rotation, and addition and subtraction modulo $2^{32}$).

## 3. Statement Regarding Trap-Doors

None known.  There are several reasons to feel confident that there are no trap-doors in this cipher.

- CAST-256 uses the four round function s-boxes in CAST-128.  The design criteria and construction procedure for these s-boxes have been published [A97, MA96] and the specific s-boxes themselves have been examined by a number of researchers.
- CAST-256 uses the three round functions in CAST-128.  The design criteria for these round functions have been published [A97] and the specific round functions themselves have been examined by a number of researchers.  Furthermore, the complexity introduced by the mixed operations in the round functions would seem to make it difficult to insert a trap-door of any kind.
- CAST-256 uses 48 rounds.  Inserting a non-obvious trap-door that will carry through 48 rounds of the cipher would seem to be a formidable task.
- CAST-256 uses a significantly more complex key scheduling algorithm than DES.  A trap-door in the final round that allows the attacker (i.e., the one knowing this trap-door) to recover information about the final round key will be of little help in deriving either other round keys or the primary key.  This contrasts with DES in which knowledge of any round key gives knowledge of the primary key with only a brute-force search over 8 bits of key.

## 4. Publications Discussing or Analyzing Aspects of the CAST Design Procedure

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[Please note that a number of the above papers are available at the following location:
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